

## Lecture 12 – February 16

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**Warning:** these notes may contain factual errors**Outline:**

- Uniform Laws via Entropy numebrs
- Classes with finite entropies
  - Nonparametric classes
  - VC classes

**Recap:** Given  $\mathcal{F}$  with distance  $d$ ,  $N(\mathcal{F}, d, \epsilon) = \min\{N \in \mathbb{N} \mid \exists \epsilon\text{-cover of } \mathcal{F} : \{f_i\}_{i=1}^N \text{ in distance } d\}$

Chaining: If  $\{X_t\}$  is sub-Gaussian process,  $\log(\mathbb{E} \exp(\lambda(X_s - X_t))) \leq \frac{\lambda^2 d(s,t)^2}{2}$ . Then,  $\mathbb{E}(\sup_{t \in T} X_t) \leq$

$$C\mathcal{J}(\mathcal{F}, d) = C \int_0^\infty \sqrt{\log N(\mathcal{F}, d, \epsilon)} d\epsilon.$$

Entropy number  $\rightarrow$  uniform laws. For empirical distirbution  $P_n$ , let  $L_p(P_n)$  be the  $L_p$  norm w.r.t.  $P_n$ , i.e.,  $\|f\|_{L_p(P_n)} = (\frac{1}{n} \sum |f(X_i)|^p)^{1/p}$ .

**Example 1:** Often use  $L_2(P_n)$  norm in symmetrized processes, i.e., if  $Z_f := \frac{1}{\sqrt{n}} \sum \epsilon_i f(X_i)$  where  $\epsilon_i \stackrel{i.i.d}{\sim} \text{unif}(-1, +1)$ .

For fixed  $X_1, \dots, X_n$ ,

$$\begin{aligned} \log(\mathbb{E} \exp(\lambda(Z_f - Z_g))) &= \log(\mathbb{E} \exp(\lambda \frac{1}{\sqrt{n}} \sum (\epsilon_i (f(X_i) - g(X_i)))) \\ &\leq \frac{\lambda^2}{2n} \sum (f(X_i) - g(X_i))^2 = \frac{\lambda^2}{2} \|f - g\|_{L_2(P_n)}^2. \end{aligned}$$

i.e.,  $f \rightarrow \frac{1}{\sqrt{n}} \sum \epsilon_i f(X_i)$  is an  $\|\cdot\|_{L_2(P_n)}^2$  sub-Gaussian process. so,

$$\mathbb{E}[\sup_{f \in \mathcal{F}} |\frac{1}{\sqrt{n}} \sum \epsilon_i f(X_i)| \mid X_1, \dots, X_n] \leq C \int_0^\infty \sqrt{\log N(\mathcal{F}, L_2(P_n), \epsilon)} d\epsilon.$$

For  $M < \infty$ , let  $f_M(x) = f(x)$  if  $|f(x)| \leq M$  or 0 otherwise. Let  $\mathcal{F}$  be a collection of functions on  $\mathcal{X}$  with envelop  $F$ , i.e.,  $|f(x)| \leq F(x)$  for all  $x \in \mathcal{X}$  and  $F \in L_1(P)$ . Define  $\mathcal{F}_M := \{f_M\}_{f \in \mathcal{F}}$ .

**Theorem 1.** (*ULLNs with entropies*): If  $\sqrt{\log N(\mathcal{F}, L_1(P_n), \epsilon)} = o_p(n)$  for all  $M < \infty, \epsilon > 0$ , then  $\|P_n - P\|_{\mathcal{F}} \xrightarrow{P} 0$ , ie.  $\mathcal{F}$  is G.C. class.

**Proof** Let  $P_n^0(f) := \frac{1}{n} \sum \epsilon_i f(X_i)$  where  $\epsilon_i \stackrel{i.i.d}{\sim} \text{unif}(-1, +1)$ . Then  $|P_n^0 f| \leq P_n |f| = \|f\|_{L_1(P_n)}$ .

$$\mathbb{E}[\|P_n - P\|] \leq 2\mathbb{E}[\sup_{f \in \mathcal{F}} |P_n^0 f|] \leq 2\mathbb{E}[\sup_{f \in \mathcal{F}} |P_n^0(f - f_M)|] + \mathbb{E}[\sup_{f \in \mathcal{F}_M} |P_n^0 f|]$$

Note that,  $2\mathbb{E}[\sup_{f \in \mathcal{F}} |P_n^0(f - f_M)|] \leq 2\mathbb{E}[F1(F \geq M)]$ .

Now control  $E[\sup_f |P_n^0 f|]$ . Let  $\mathcal{G}$  be a minimal  $\epsilon$  - cover of  $\mathcal{F}_M$  in  $L_1(P_n)$  norm. Then,  $\text{card}(\mathcal{G}) = N(\mathcal{F}_M, L_1(P_n), \epsilon)$ .

Therefore,  $\sup_{f \in \mathcal{F}} |P_n^0 f| \leq \max_{g \in \mathcal{G}} |P_n^0 g| + \epsilon$  (by triangular inequality).

Note that w.l.o.g.  $|g(x)| \leq M$ , all  $g \in \mathcal{G}$ , so  $P_n^0 g$  is  $\frac{M^2}{n}$  sub- Gaussian. So

$$\mathbb{E}[\max_{g \in \mathcal{G}} |P_n^0 g| \mid X] \leq 2\sqrt{2\frac{M^2}{n} \log N(\mathcal{F}, L_1(P_n), \epsilon)}$$

LHS is less than  $M$ , and the RHS is  $\sqrt{\frac{M}{n} o_p(n)} = o_p(1)$

$$\begin{aligned} \mathbb{E}[\mathbb{E}[\max_g |P_n^0 g|]] &\leq \mathbb{E}[\min(M, o_p(1))] \rightarrow 0 \text{ as } n \rightarrow \infty. \\ \Rightarrow \mathbb{E}[\|P_n - P\|_{\mathcal{F}}] &\leq 2\mathbb{E}[F1(F \geq M)] + o(1) + \epsilon \end{aligned}$$

Since,  $\mathbb{E}[F1(F \geq M)] \rightarrow 0$  as  $n \rightarrow \infty$ , the proof is done.  $\square$

Understand uniform entropies: Often random covering numbers such as  $N(\mathcal{F}, L_r(P_n), \epsilon)$  are a bit annoying. so try to give conditions such that  $\sup_P N(\mathcal{F}, L_r(P), \epsilon)$  can be controlled.

Let's look at some examples in non-parametric function classes.

**Example 2:** Let  $\mathcal{F}$  be the collection of 1- Lipschitz functions on  $[0, 1]$  with  $f(0) = 0$ . Fix  $\epsilon > 0$ , consider  $\|f\|_{\infty} := \sup_{x \in [0, 1]} |f(x)|$ . By dividing the unit intervals by intervals with length  $\epsilon$  and moving along x axis by epsilon with 3 choice of directions, namely up(45 degree angle), staright, down(45 degree angle) Packing #s  $\geq 3^{\frac{1}{\epsilon}}$  (1-Lipschitz function's height change associated with change of  $\epsilon$  is also at most  $\epsilon$ ). Thus,  $\sup_{P: \text{supp } P = [0, 1]} \log N(\mathcal{F}, L_r(P), \epsilon) \leq \frac{C}{\epsilon}$  where  $c < \infty$  is absolute constant. So

$$\mathbb{E}[\sup |P_n - Pf|] \leq 2\mathbb{E}[\sup_f |P_n^0 f|] \leq \frac{c}{\sqrt{n}} \mathbb{E}[\int_0^1 \sqrt{\log N(\mathcal{F}, L_2(P_n), \epsilon)} d\epsilon] \leq \frac{c}{\sqrt{n}} \int_0^1 \frac{1}{\sqrt{\epsilon}} d\epsilon \leq \frac{c}{\sqrt{n}}$$

In 2+ dimensions, divided boxes with length  $\epsilon$  has  $\frac{1}{\epsilon^2}$  boxes, (or  $(\frac{1}{\epsilon})^d$  in  $d$  dimensions), so

$$\log N(\mathcal{F}, \|\cdot\|_{\infty}, \epsilon) \geq \frac{c}{\epsilon^2} \Rightarrow \mathcal{J}(\mathcal{F}, \|\cdot\|_{\infty}) = \int_0^1 \frac{1}{\epsilon} = +\infty.$$

Vapnik- Chervonenkis (VC classes) Collections of functions or sets with nice combinatorial structure allowing uniform entropy/covering number bounds.

**Definition 0.1.** Let  $\mathcal{C}$  be a collection of sets and  $X = \{X_1, \dots, X_n\}$  be a collection of points . A vector  $y \in \{+1, -1\}^n$  is a labeling of  $X$ . Say  $\mathcal{C}$  shatters  $X$  if for all labelings  $y$  of  $X$ ,  $\exists$  a set  $A \in \mathcal{C}$ , i.e.,  $X_i \in A$  if  $y_i = 1$  and  $X_i \notin A$  if  $y_i = -1$ .

Equivalently,  $\{x_1, \dots, x_n\} \cap \mathcal{C} = \{A \cap \{x_1, \dots, x_n \mid A \in \mathcal{C}\} = 2^X$ .

**Example 3:** Let  $x_1, x_2, x_3 \in \mathbb{R}^2$ , not collinear.  $\mathcal{C} = \{\text{half space in } \mathbb{R}^2\}$ . Then  $\mathcal{C}$  shatters  $\{x_1, x_2, x_3\}$

**Definition 0.2.** : The VC- dimension  $VC(\mathcal{C})$  is the size of the largest set  $\{x_1, \dots, x_n\}$  s.t.  $\mathcal{C}$  shatters  $\{x_1, \dots, x_n\}$ .

**Definition 0.3.**  $\Delta_n(\mathcal{C}, \{x_1, \dots, x_n\}) :=$  the number of labelings  $\mathcal{C}$  realizes on  $\{x_i\}$ . Then  $VC(\mathcal{C}) := \sup\{n \in \mathbb{N} \mid \max_{x_1, \dots, x_n} \Delta_n(\mathcal{C}, \{x_i\}) = 2^n\}$ .

**Example 4:** Half-spaces in  $\mathbb{R}^d$  have  $VC(\mathcal{C}) = d + 1$ , Think of  $\mathbb{R}^2$ . Then  $VC(\mathcal{C}) \geq 3$ . To do rigorously requires arguing (by geometry) that we would have to have the situation where diagonal labeling does not work.

**Lemma 2.** (Sauer- Shelah) for any class  $\mathcal{C}$ ,

$$\max_{x_1, \dots, x_n} \Delta_n(\mathcal{C}, \{x_i\}) \leq \sum_{k=0}^{VC(\mathcal{C})} \binom{n}{k} = O(n^{VC(\mathcal{C})}).$$

Consequence: If  $\sup_{x_1, \dots, x_n} \Delta_n(\mathcal{C}, \{x_i\}) < 2^n$ , then  $\Delta_n(\mathcal{C}, \{x_i\})$  is polynomial in  $n$ .

Let  $L_r(P)$  norm on sets  $A \subset \mathcal{X}$  be defined by  $\|1_A\|_{L_r(P)} = (\int 1(x \in A)^r dP(x))^{1/r}$

**Theorem 3.** :  $\exists$  a universal constant  $K < \infty$  s.t.  $\forall \epsilon > 0$ ,

$$\sup_P N(\mathcal{C}, L_r(P), \epsilon) \leq K \cdot VC(\mathcal{C}) \cdot (4e)^{VC(\mathcal{C})} \left(\frac{1}{\epsilon}\right)^{r \cdot VC(\mathcal{C})}$$

$$\Rightarrow \log N(\mathcal{C}, L_r(P), \epsilon) \leq c \cdot r \cdot VC(\mathcal{C}) \cdot \log(1/\epsilon)$$